Recognizing String Graphs in NP

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Identification des graphes de corde dans NP

(Recognizing String Graphs in NP)

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The origin of the problem

Sinden 1966 Topology of Thin Film RC Circuits



	1	2	3	4
1		X		X
2	X		X	X
3		X		X
4	X	X	X	





Is G an intersection graph of a set of curves in the plane?

	1	2	3	4
1		X		X
2	X		X	X
3		X		X
4	X	X	X	

String Graphs



Planar graphs are string graphs (Sinden, 1966)

Recognizing string graphs is NP-hard (Kratochvíl, 1991)

Recognizing string graphs is decidable (in NEXP) (Pach, Tóth, 2000; Schaefer, Š, 2000)



Can G be drawn in the plane ?

red edge may intersect green edge
red edge may intersect orange edge
no other pair of edges may intersect

String Weak realizability (Matoušek, Nešetřil, Thomas'88)

Weak realizability

Input:Graph Gset R of pairs of edges

Output: Is there a drawing of G in the plane such that only pairs from R may intersect?

(e.g. R=0 corresponds to planarity testing)

Weak realizability

NP-hard (Kratochvíl '91) NEXP (Pach, Tóth '00; Schaefer, Š '00)

Theorem: If there is a drawing realizing (G,R) then there is a drawing with at most m2^m intersections where m is the number of edges of G.

The Theorem is tight (Kratochvíl, Matoušek '91)



edge properly embedded arc (parc)

isotopy rel endpoints = continuous
deformations not moving endpoins

Intersection number i(α,β) of two parcs α,β min{|a b|; a C(α), b C(β)}

(set of curves isotopic to α)

Lemma: On an orientable surface any collection of parcs can be redrawn so that any two parcs α , β intersect at most i(α , β) times.



encode the properly embedded arcs (up to isotopy)
for each pair α,β not in R check i(α,β)=0

Encoding the parcs 1) A triangulation T of M



Encoding the parcs 2) Normalization of the parc w.r.t. T



Encoding the parcs 3) Compute normal coordinates



Encoding the parcs

Parcs having the same normal coordinates are isotopic rel boundary.





Encoding the parcs Is it polynomial ?

construct a weak realizability problem including the triangulation and use

Theorem: If there is a drawing realizing (G,R) then there is a drawing with at most m2^m intersections where m is the number of edges of G.

Word equations x,y – variables xayxb=axbxy a,b - constants a solution x=aaaa v=b Word equations with given lengths xayxb=axbxy The size of the bit **X** =4 representation of the numbers counts V =1 to the size of the input **Word equations**

NP-hard in PSPACE (Plandowski '99)

Word equations with given lengths

in P (Plandowski, Rytter '98) the lexicographically smallest solution given by a straight line program

Coloring components of a curve

normal coordinates – can encode any embedded collection of closed curves and parcs (=curve)



x+y=a x+z=b y+z=c



The proof of weak realizability

- encode the properly embedded arcs (up to isotopy)
 for each pair α,β not in R check i(α,β)=0
 - Do coordinates of α encode a parc? Are parcs α , β isotopically disjoint?
 - check that both α , β are parcs
 - color one component of α + β by "b"
 - They are disjoint iff the component is either α or β

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Consequences + other results pairwise crossing number NP existential theory of diagrams (topological inference) NP



A intersects B B intersects C A is disjoint from C

weak realizability on different surfaces Can be done in NP?